

Discrete Structures

Fight for Survival

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August 27, 2026



Quantifiers

A predicate can become a proposition when we assign it a fixed value from the universe of discourse.

But there's another powerful tool: **quantification**. This allows us to determine whether the predicate is true (or false) for all possible values in the universe of discourse or for some specific value(s).

- **Universal Quantifier ($\forall$)**: True or false for **all** values in the universe of discourse.
- **Existential Quantifier ($\exists$)**: True or false for **some** values in the universe of discourse.



Universal Quantification: Examples

Universal quantification ($\forall$) is like a magic spell that makes statements apply to **every** member of a set.

Let's see some fun examples:

- $\forall$ **Cats**: Love napping in sunbeams.
- $\forall$ **Programmers**: Drink coffee by the gallon.
- $\forall$ **Professors**: Claim their subject is the most interesting.

Universal Quantifier - Example

Let $P(x)$ = "x must take the Discrete Structure course",
and $Q(x)$ = "x is a Computer Science student".

The universe of discourse is "all students of ITU (x)".

- Express the statement in symbolic form: "All computer science students must take the Discrete Structure course"

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$$\forall x(Q(x) \rightarrow P(x))$$

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$$\forall x (Q(x) \vee P(x))$$

Existential Quantifier

Definition

The existential quantification of a predicate $P(x)$ is the proposition " $P(x)$ is true for some of the values of x in the universe of discourse". We use the notation

$$\exists x P(x)$$

which is read as "there exists x ".

Similarly, if the universe of discourse is finite, say $\{n_1, n_2, n_3, \dots, n_k\}$, then the existential quantifier is simply the disjunction of all the elements

$$\exists x P(x) \equiv P(n_1) \vee P(n_2) \vee P(n_3) \vee \dots \vee P(n_k)$$



Existential Quantification: Examples

Existential quantification ($\exists$) allows us to find **at least one** instance where a statement holds true.

Let's explore some practical examples:

- $\exists$ **Students Who Love Math**: In a class of 60, some are passionate about math.
- $\exists$ **Available Parking Spots**: In a busy Arfa Tower, there are spots waiting to be found.

Truth values of Quantifiers

In general, when are quantified statements **true/false**?

Statement	True When	False When
$\forall x P(x)$	$P(x)$ is true for every x	There is an x for which $P(x)$ is false
$\exists x P(x)$	There is an x for which $P(x)$ is true	$P(x)$ is false for every x

Table: Truth Conditions of Universal and Existential Quantifiers

Example - Express in Mathematical Notation

- The product of two negative integers is positive.
- The difference of two negative integers is not necessarily negative.

Example - Express in Mathematical Notation

- The product of two negative integers is positive.

$$\forall x, y \in \mathbb{Z}^-, x \times y > 0$$

- The difference of two negative integers is not necessarily negative.

$$\exists a, b \in \mathbb{Z}^-, a - b > 0$$

Precedence of Quantifiers

The quantifiers $\forall$ and $\exists$ have higher precedence than all the other logical operators.

For example,

$$\exists a P(a) \vee Q(a) \equiv (\exists a P(a)) \vee Q(a)$$

$$\exists a P(a) \vee Q(a) \not\equiv \exists a (P(a) \vee Q(a))$$



Mindful of the Domain

Always be **mindful of the domain** when dealing with quantifiers. The same statement can be true or false based on the domain.

For example:

- $\forall x(x^2 > x)$
 - Domain: x is a positive integer
 - **This statement is false.**
 - Counterexample: $x = 1$
- $\forall x(x^2 > x)$
 - Domain: x is a positive integer greater than 1
 - **This statement is true.**

Proving/Disproving Quantifiers

To prove **universally** quantified statements, we need to show that the statement is true for **all possible values** in the domain.

Proving/Disproving Quantifiers

To disprove **universally** quantified statements, we need to show that the statement is false for **just one value** in the domain. (also referred to as the **counterexample**).

Proving/Disproving Quantifiers

To prove **existentially** quantified statements, we need to show that the statement is true for **just one value** in the domain.

Proving/Disproving Quantifiers

To disprove **existentially** quantified statements, we need to show that the statement is false for **every value** in the domain.

A question:

If a universally quantified statement is true, does that mean the existentially quantified statement is always true?

Free and Bound variables

- A variable x in the predicate $P(x)$ is called a **free variable** because the variable is **free** to take on any value in the domain.
- The variable x in the statement $\forall xP(x)$ or $\exists xP(x)$ is a **bound variable** because **the variable is bound** to the quantifier within the domain.
- A **statement with no free variables** is a proposition because the statement's truth value can be determined.
- The part of **logical expression** to which a quantifier is applied is called the **scope** of the quantifier.

Free and Bound variables

- $\forall x P(x) \wedge Q(x)$ is **not a proposition**.
- **But** $\forall x (P(x) \wedge Q(x))$ is **a proposition**.

Quantifiers - Negation (DeMorgan's Law)

$\forall x P(x)$: Every student in this class has taken a calculus course.



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What is the negation of the above statement, that is, $\neg(\forall xP(x))$?

It is not true that every student in this class has taken calculus.

There is at least one student that has not taken calculus.

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$$\neg(\forall xP(x)) \equiv \exists x\neg P(x)$$

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Domain of discourse $\{x_1, x_2, \dots, x_n\}$

$$\neg \forall x P(x)$$

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$$\neg \forall x P(x)$$



$$\neg (P(x_1) \wedge P(x_2) \wedge P(x_3) \wedge \dots \wedge P(x_n))$$

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$$\equiv \neg P(x_1) \vee \neg P(x_2) \vee \neg P(x_3) \vee \dots \vee \neg P(x_n)$$

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$$\neg \forall x P(x) \equiv \neg (P(x_1) \wedge P(x_2) \wedge P(x_3) \wedge \dots \wedge P(x_n))$$

$\equiv$

$$\exists x \neg P(x) \equiv \neg P(x_1) \vee \neg P(x_2) \vee \neg P(x_3) \vee \dots \vee \neg P(x_n)$$

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$$\forall x \neg P(x) \equiv \neg P(x_1) \wedge \neg P(x_2) \wedge \neg P(x_3) \wedge \dots \wedge \neg P(x_n)$$