

# Discrete Structures

## Fight for Survival

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August 27, 2026



## Solving Recurrences

# Recurrences

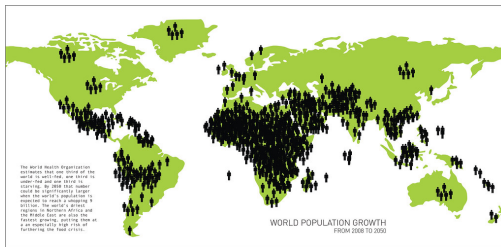
We observe that things are highly correlated to their immediate past!

# Recurrences



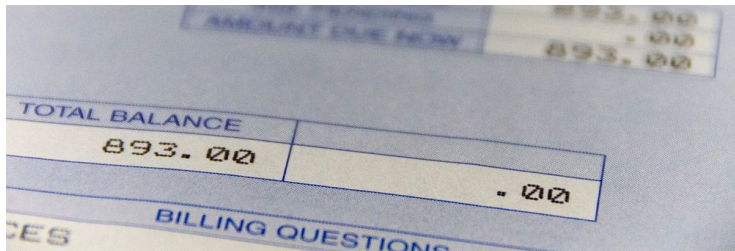
Temperature now is very close to what it was one minute ago!

# Recurrences



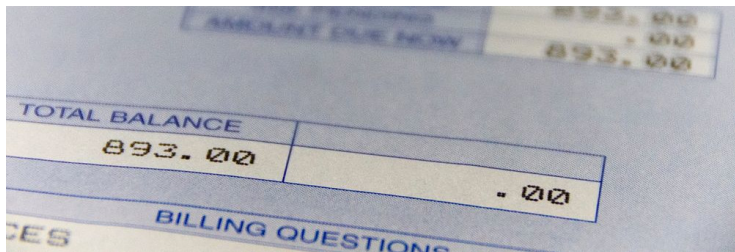
World Population is very reliable function of what it was a year ago!

# Recurrences



Your bank account balance looks very much like what it was a month ago!

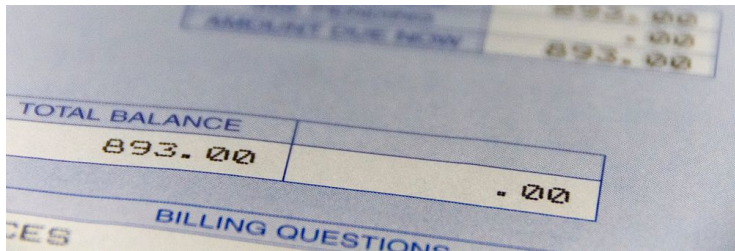
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We can reliably estimate current state of things as a function of previous measurements. This is called a recurrence.

Logistic map, factorials, fibonacci sequences etc.

# Solving Recurrences

- Many famous/important algorithms use the **recursion**.
- For example, **MergeSort**, **QuickSort**, **Tree searching**, **Strassen's algorithm** etc.
- Time Complexity of these algorithms is written as a **recurrence**.
- There are multiple ways to solve a **recurrence**.
- We will see a couple of them today.

# Solving Recurrences

- Many famous/important algorithms use the **recursion**.
- For example, **MergeSort**, **QuickSort**, **Tree searching**, **Strassen's algorithm** etc.
- Time Complexity of these algorithms is written as a **recurrence**.
- There are multiple ways to solve a **recurrence**.
- We will see a couple of them today.

But we need some practice on **summation of series** and finding **closed forms**.

# Summation and closed form

A **closed form** is an expression that involves very few operations ( $+$ ,  $-$ ,  $\times$ ,  $\div$  etc).

A **closed form** is an expression that can be computed by applying a fixed number of familiar operations to arguments (e.g.,  $n(n+1)$  ).

Determining the closed form can make it easier to interpret.  
The following expression:  $2 + 4 + \dots + 2n$  (**is not closed form**).

# Summation and closed form

Some helpful facts for deriving closed forms:

- $\sum_{k=1}^n c = c(n)$  [Sum of Constant]
- $\sum_{k=m}^n c = c(n - m + 1)$  [Sum of Constant]
- $\sum ca_k = c \sum a_k$  [Sum of Constant]
- $\sum_{k=1}^n (a_k - a_{k-1}) = a_n - a_0$  [Collapsing sum]
- $\sum_{k=1}^n (a_{k-1} - a_k) = a_0 - a_n$  [Collapsing sum]
- $\sum (a_k + b_k) = \sum (a_k) + \sum (b_k)$  [Sum of sums]
- $\sum_{k=m}^n a_{k+i} = \sum_{k=m+i}^{n+i} a_k$  [Shifting index]

# Summation and closed form

Some well-known closed forms that everyone should know:

- $\sum (a_k x^{i+k}) = x^i \sum (a_k x^k)$
- $\sum_1^n k = \frac{(n(n+1))}{2}$
- $\sum_{k=1}^n k^2 = \frac{(n(n+1)(2n+1))}{6}$
- $\sum_{k=0}^n a^k = \frac{a^{n+1}-1}{a-1} \quad a \neq 1$
- $\sum_{k=1}^n k a^k = \frac{a-(n+1)a^{n+1}+na^{n+2}}{(a-1)^2} \quad a \neq 1$

# Summation and closed form

Let's take a look at how we might convert a given summation into a closed form:

## Example

Find a closed form for the following:

$$3 + 7 + 11 + \dots + (3 + 4n)$$

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## Example

Find a closed form for the following:

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**Step 1:** Get a handle on the summation

$$= 3 + 7 + 11 + \dots + (3 + 4n) = \sum_{k=0}^n (3 + 4k)$$

**Step 2:** Try to simplify or apply any known tricks

$$= 3(n+1) + \sum_{k=0}^n (4k) = 3(n+1) + \sum_{k=1}^n (4k)$$

$$= 3(n+1) + 4 \sum_{k=1}^n k = 3(n+1) + 4\left(\frac{n(n+1)}{2}\right)$$

$$= 3 + 3n + 2n^2 + 2n = 2n^2 + 5n + 3$$

# Summation and Closed forms

Try the following.

## Example

Find a closed form for the following expression:

$$\sum_{k=1}^n (2k + 3)$$

# Summation and Closed forms

Try the following.

## Example

Find a closed form for the following expression:

$$\sum_{k=1}^n (2k + 3)$$

$$\begin{aligned}\sum_{k=1}^n (2k + 3) &= 3n + 2 \sum_{k=1}^n k \\ &= 3n + 2\left(\frac{n(n+1)}{2}\right) = 3n + n(n+1) = n^2 + 4n\end{aligned}$$

# Summation and Closed forms

Try the following.

Example

Find a closed form for the following sequence:

$$4 + 8 + 12 + 16 + \dots + 4n$$

# Summation and Closed forms

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## Example

Find a closed form for the following sequence:

$$4 + 8 + 12 + 16 + \dots + 4n$$

$$4i$$

[Try to find the pattern (find the  $i^{th}$  term)]

$$\sum_{i=1}^n 4i$$

[Write the summation expression]

$$= \sum_{i=1}^n 4i$$

[Simplify- find the closed form]

$$= 4 \sum_{i=1}^n i$$

$$= 4 \left( \frac{n(n+1)}{2} \right)$$

$$= 2n^2 + 2n$$



# Summation and Closed forms

Try the following.

## Example

Find a closed form for the following (a bit challenging):

$$2 + (2^2 \times 7) + (2^3 \times 14) + (2^4 \times 21) + (2^n \times 7(n-1))$$

## Answer

$$2^i \times 7(i-1)$$

[Try to find the pattern (find the  $i^{th}$  term)]

$$= 2 + \sum_{i=1}^n 2^i \times 7(i-1)$$

[Write the summation expression]

$$= 2 + 7 \sum_{i=1}^n 2^i (i-1)$$

[Simplify, take constant out (Fact 2)]

$$= 2 + 7 \sum_{i=1}^n (2^i i - 2^i)$$

[Simplify, (Fact 5)]

$$= 2 + 7 \left( \sum_{i=1}^n 2^i i - \sum_{i=1}^n 2^i \right)$$



# Summation and Closed forms

## Example

Find a closed form for the following (a bit challenging):

$$2 + (2^2 \times 7) + (2^3 \times 14) + (2^4 \times 21) + (2^n \times 7(n-1))$$

$$= 2 + 7(\sum_{i=1}^n 2^i i - \sum_{i=1}^n 2^i)$$

$$= 2 + 7((2 - (n+1)2^{n+1} + n2^{n+2}) - \sum_{i=1}^n 2^i)$$

$$= 2 + 7((2 - (n+1)2^{n+1} + n2^{n+2}) - (2^{n+1} - 1 - 1))$$

$$= 2 + 7((2 + (n-1)2^{n+1}) - (2^{n+1} - 1 - 1))$$

$$= 2 + 7((2 + (n-1)2^{n+1}) - (2^{n+1} - 2))$$

$$= 2 + 7(4 + (n-2)2^{n+1})$$

[Simplify, (Fact 11)]

[Simplify, (Fact 10)]

[Simplify red term]

[Simplify blue term]

[Further  
simplification]



# Summation and closed forms

## Example

Let  $\text{count}(n)$  be the number of  $:=$  statements executed by the following algorithm as a function of  $n$ . Find a closed form for a  $\text{count}(n)$ .

```
i := 1;
while i < n
  i := i + 1;
  for j := 1 to i
    do S
  end
end
```

1 assignment op  
 $(n - 1)$

$\times (2 + 3 + \dots + n)$  ops

# Summation and closed forms

Answer:

The total count of operations,  $\text{count}(n)$  is:

$$\begin{aligned} &= 1 + (n - 1) + (2 + 3 + 4 + \dots + n) \\ &= 1 + 2 + 3 + 4 + \dots + n + (n - 1) = \frac{n(n+1)}{2} + (n - 1) \\ &= \frac{n(n+1)}{2} + \frac{2(n-1)}{2} = \frac{n^2+3n-2}{2} \end{aligned}$$

# Summation and closed forms

Lets try a harder one.

## Example

Let  $\text{count}(n)$  be the number of  $:=$  statements executed by the following algorithm as a function of  $n$ . Find a closed form for a  $\text{count}(n)$ .

```
i := 1;
while i < n
  i := i + 2;
  for j := 1 to i
    do S
  end
end
```

1  
 $(n-1)/2$

(i goes up by 2) ops  
???

# Summation and closed forms

Lets try a harder one.

Let's pick a specific  $n$  (i.e.,  $n = 8$ ) to just try to get a handle on what the heck is happening inside the loop.

```
i := 1;  
while i < n  
  i := i + 2;  
  for j := 1 to i  
    do S  
  end  
end
```

| i | i < 8? | j      |
|---|--------|--------|
| 1 | T      | 1 to 3 |
| 3 | T      | 1 to 5 |
| 5 | T      | 1 to 7 |
| 7 | T      | 1 to 9 |
| 9 | F      |        |

# Summation and closed forms

Lets try a harder one.

Let's pick a specific  $n$  (i.e.,  $n = 8$ ) to just try to get a handle on what the heck is happening inside the loop. Assume  $n = 2k$ .

```
i := 1;
while i < n
  i := i + 2;
  for j := 1 to i
    do S
  end
end
```

| i    | i < n? | j         | calls to S |
|------|--------|-----------|------------|
| 1    | T      | 1 to 3    | 3          |
| 3    | T      | 1 to 5    | 5          |
| 5    | T      | 1 to 7    | 7          |
| 7    | T      | 1 to 9    | 9          |
| 2k-1 | T      | 1 to 2k+1 | 2k+1       |
| 2k+1 | F      |           |            |

# Summation and closed forms

So the number of calls to  $S$  looks like  
calls to  $S$ :

$$\begin{aligned} & 3 + 5 + 7 + \dots + (2k + 1) \\ &= \sum_{h=1}^k (2h + 1) = \sum_{h=1}^k (2h) + \sum_{h=1}^k (1) \\ &= 2 \times \frac{k(k+1)}{2} + k \\ &= k^2 + 2k \end{aligned}$$

We now know the total number of calls to  $S$ . But what is  $k^2 + 2k$  in terms of  $n$ ?

# Summation and closed forms

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We now know the total number of calls to  $S$ . But what is  $k^2 + 2k$  in terms of  $n$ ?

$$k^2 + 2k = n/2(n/2 + 2) = \Theta(n^2)$$



# Summations and Recurrence Relation

Next :

What if sequence is given in the form of a recurrence relation (instead of individual terms) ?

$$a(n) = c_1 a(n-1) + c_2$$

How can we find the closed form of the recurrence relation  $a(n)$  ?

# Recurrence Relations : Substitution method

Given a recurrence relation:

$$a(n) = a(n-1) + 2n; \quad a(0) = 1$$

Find the closed form of  $a(n)$ .

**Step1:**

Work **backwards to unravel** and write out some terms.

$$a(n) = a(n-1) + 2n$$

$$a(n-1) = a(n-2) + 2(n-1)$$

$$a(n-2) = a(n-3) + 2(n-2)$$

**Step2:**

Use **substitutions** and plug in for the terms.

$$a(n) = a(n-1) + 2n$$

$$a(n) = a(n-2) + 2(n-1) + 2n$$

$$a(n) = a(n-3) + 2(n-2) + 2(n-1) + 2n$$

# Recurrence Relations : Substitution method

Given a recurrence relation:  $a(n) = a(n-1) + 2n$ ;  $a(0) = 1$  Find the closed form of  $a(n)$ .

## Step3:

Look for a **pattern** among the terms and visualize the terms over its entirety.

$$a(n) = a(n-3) + 2(n-2) + 2(n-1) + 2n$$

$$i^{th} \text{ term} = 2i$$

## Step4:

Establish the **summation formula** and then determine the closed form

$$\sum_{i=1}^n (2i) = 2\left(\frac{n(n+1)}{2}\right) = n^2 + n$$

So, the final closed form would be:  $n^2 + n + 1 = O(n^2)$  add 1 for initial term  $a(0)$



# Substitution method-More examples

**Practice:** Analyze how much work a recursive algorithm does by precisely calculating its closed form based on the following recurrence relation.

$$r(n) = 2r(n - 1) + 1; \quad r(0) = 1$$

# Substitution method-More examples

**Step 1:** Work backwards and write out some terms:

$$r(n) = 2r(n-1) + 1$$

$$r(n-1) = 2r(n-2) + 1$$

$$r(n-2) = 2r(n-3) + 1$$

$$r(n-3) = 2r(n-4) + 1$$

**Step 2:** Use substitutions for the terms...

$$r(n) = 2r(n-1) + 1$$

$$= 2[2r(n-2) + 1] + 1$$

$$= 2^2 r(n-2) + 2 + 1$$

$$= 2^2 [2r(n-3) + 1] + 2 + 1$$

$$= 2^3 r(n-3) + 2^2 + 2 + 1$$

Do you see a pattern? This looks like a geometric sequence (there's a summation formula for that).

$$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^n$$



## Substitution method-Cont...

**Step 4:** Try to establish the summation formula and determine the closed form.

$$\sum_{k=0}^n a^k = \frac{a^{n+1}-1}{a-1}; a \neq 1$$

This leads us to a closed form of  $2^{n+1} - 1$  which is  $O(2^n)$ .

# Recurrence Relations: Cancellation Method

Another technique that is useful for solving a recurrence relation is the:

Cancellation method.

# Recurrence Relations: Cancellation Method

How it works.

- 1 Start with the general (recurrence) equation  $r(n)$  and continue writing new equations whose left side is the term involving the recurrence  $r$  on the right side of the **previous equation**.
- 2 Look for a pattern that allows us to skip ahead and write the last equation  $r(0)$  on the right side.
- 3 Add up the equations and cancel the like terms to obtain an expression for  $r(n)$ .

# Recurrence Relations: Cancellation Method

**Example** Solve the following recurrence using cancellation:

$$a(0) = 1$$

$$a(n) = a(n-1) + 2n$$

**Step1:** As with substitution, we start by writing out some terms:

$$a(n) = a(n-1) + 2n$$

$$a(n-1) = a(n-2) + 2(n-1)$$

$$a(n-2) = a(n-3) + 2(n-2)$$

....

$$a(n - (n-1)) = a(0) + 2(n - (n-1)) \text{ same as } a(1) = a(0) + 2$$

# Recurrence Relations: Cancellation Method

**Step 2:** Next we add up the terms on both sides of the = sign

$$a(n) + \cancel{a(n-1)} + \cancel{a(n-2)} + \cancel{a(n-3)} + \dots + \cancel{a(1)} + \cancel{a(0)}$$
$$=$$

$$\cancel{a(n-1)} + 2n + \cancel{a(n-2)} + 2(n-1) + \cancel{a(n-3)} + 2(n-2) + \dots + \cancel{a(0)} + 2 + a(0)$$

**Step 3:** Cancel the like terms to obtain the equation for  $a(n)$ .

$$a(n) = a(0) + 2(1) + \dots + 2(n-2) + 2(n-1) + 2n$$

# Recurrence Relations: Cancellation Method

**Step 4:** Determine the summation and closed form.

$$a(n) = a(0) + 2(1) + \dots + 2(n-2) + 2(n-1) + 2n$$

$$a(n) = 1 + 2[1 + \dots + (n-2) + (n-1) + n]$$

$$a(n) = 1 + 2\left[\frac{n(n+1)}{2}\right]$$

$$a(n) = n^2 + n + 1$$

This is an  $O(n^2)$  algorithm

# Recurrence Relations: Tree Method

**Tree Method** is a very intuitive method to deduce or guess the solution to a recurrence pictorially.

We will study and see examples of this method later once we have defined graphs and trees.